
Diagrammatic Software Specification

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Introduction

- Need for a common language for
 - Designers
 - Programmers
 - Experts of the universe
 - ...
- In order to:
 - reason about the problem
 - verify business logic
- Graphical language is well suited for this task

Why a graphical language?

- Simple
 - Easier to understand compared to textual languages
- Universal
 - Different background → same understanding

What's a good modelling language?

- Graph-based,
- has a formal interpretation and
- sufficiently expressive (to capture all the peculiarities of the universe)

Proposal: GENERALIZED SKETCHES

State-of-the-art

For specifications in software engineering:

- Many graphical modeling languages
- Few of them have proper semantics

→ Specifications that are difficult to maintain due to:

- ambiguous constructions
- semantic relativism

Our contribution to the field

- Use of Generalized Sketches (GS) to:
 - compare
 - unify
 - integratediagrammatic specification languages
- Achieve this by:
 - Mapping different modeling languages to the GS format
 - Superimposing signatures from different graphical notations
- Graphical notations are parameterized by signatures
- What should be modeled must be standardized
- BNF is an analogy:
a formal notation to describe the syntax of a given language

Generalized Sketches: history

- Ehresmann's sketches: graphical presentations of categories (1968)
- Makkai: (1995) diagram predicates + sketches → Generalized Sketches.
- Diskin (1997) made GS:
 - more direct
 - more applicable to SE
 - complex diagram operations (operational ske.)

Formal definitions

- Graph: $G = (G_0, G_1, src^G, trg^G)$
- Graph homomorphism $h: G \rightarrow H$
 - $h_0: G_0 \rightarrow H_0$
 - $h_1: G_1 \rightarrow H_1$
 - $\forall e \in G_1:$
 - $src^H(h_1(e)) = h_0(src^G(e))$ and
 - $trg^H(h_1(e)) = h_0(trg^G(e))$
- Diagram: $\delta: Shape \rightarrow G$
- Signature: $\Sigma = \langle P, Arity(P) \rangle$
- Sketch: $(\Sigma -) sketch S = \langle G_S, D_S(P) \rangle$
where $\{\delta: Arity(P) \rightarrow G_S\} \in D_S(P)$

Sketch morphism

*A sketch morphism $\mu: S_1 \rightarrow S_2$
is a graph homomorphism $\mu: G_{S_1} \rightarrow G_{S_2}$
such that $\forall \delta: \text{Arity}(P) \rightarrow G_{S_1} \in D_{S_1}$
 $\Rightarrow \delta; \mu: \text{Arity}(P) \rightarrow G_{S_2} \in D_{S_2}$*

$$\Sigma_{UML} = (P , Arity(P))$$

concept

value

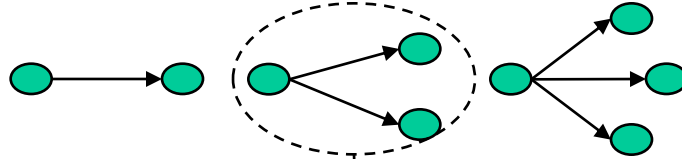
total, cover and single valued

partial and single valued

total and single valued

inclusion

[1-1]

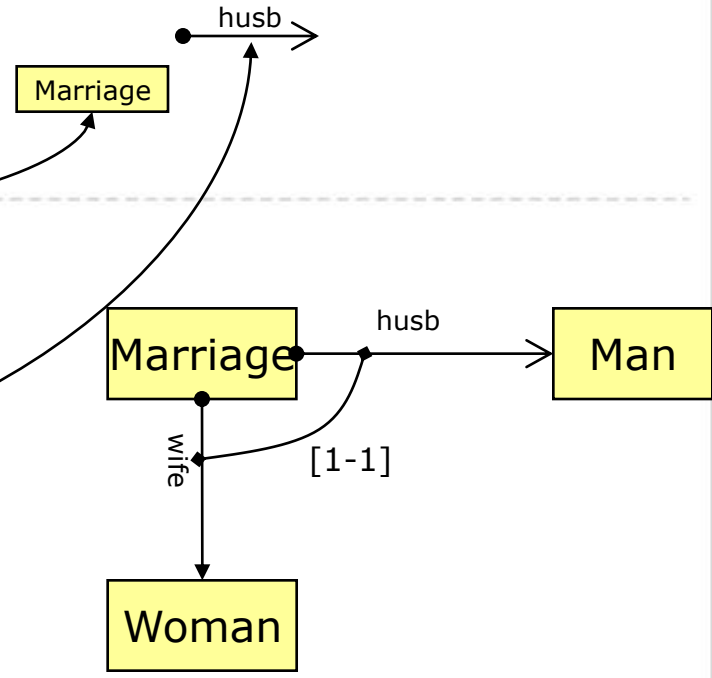
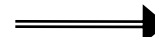
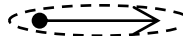
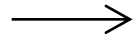
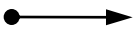


Arity(P)

Concept name

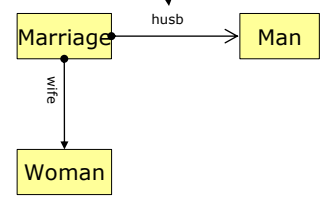
[value type]

[String], [Int]



$$S = (G_S , D_S(P))$$

G_S : carrier graph,
 $(d: Arity(P) \rightarrow G_S)$ in $D_S(P)$



$$\Sigma_{UML} = (P , \text{Arity}(P))$$

concept

Concept name

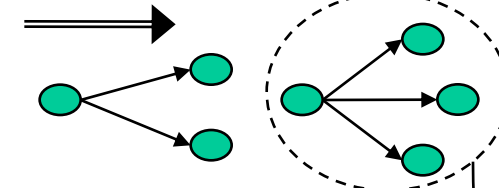
value

[value type] [String], [Int]

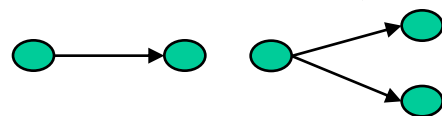
total, cover and single valued
partial and single valued
total and single valued



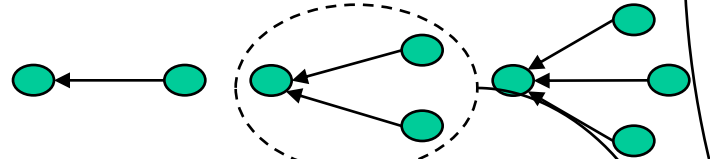
inclusion



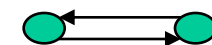
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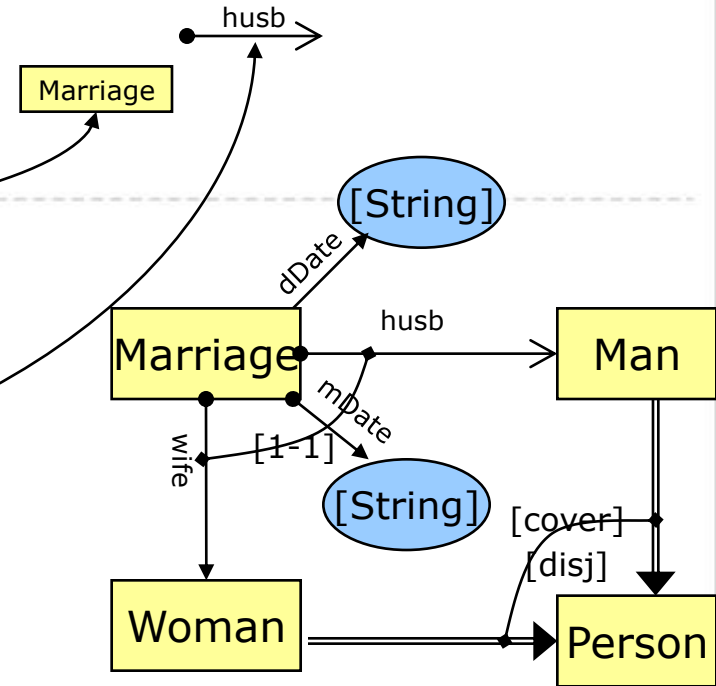
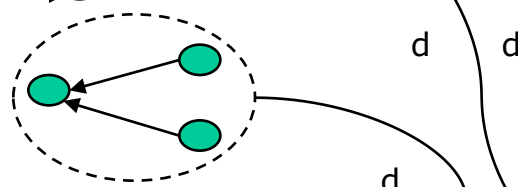
[cover]



[inv]

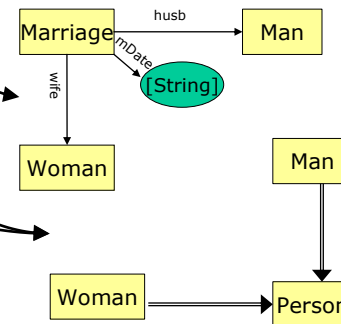


[disj]



$$S = (G_S , D_S(P))$$

G_S : carrier graph,
 $(d: \text{Arity}(P) \rightarrow G_S)$ in $D_S(P)$



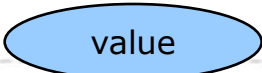
$$\Sigma_{ER} = (P, \text{Arity}(P))$$

concept

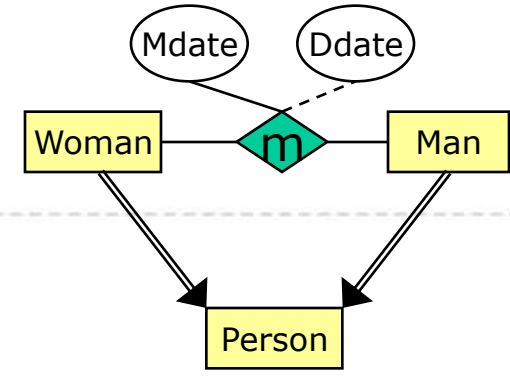
Concept name



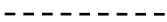
value



Mdate, Ddate



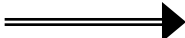
Partial mapping



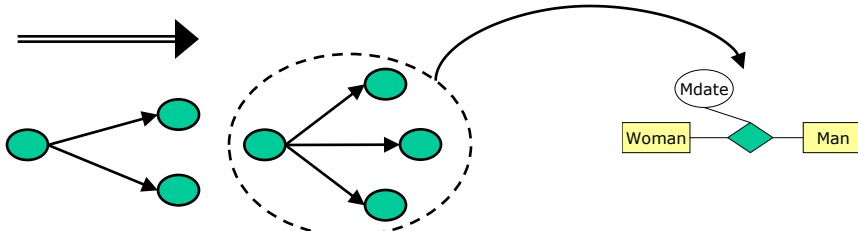
Total mapping



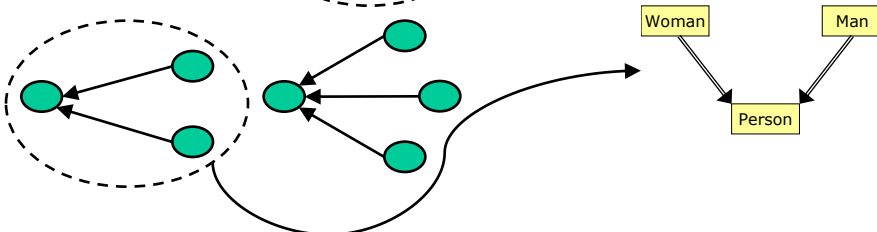
inclusion



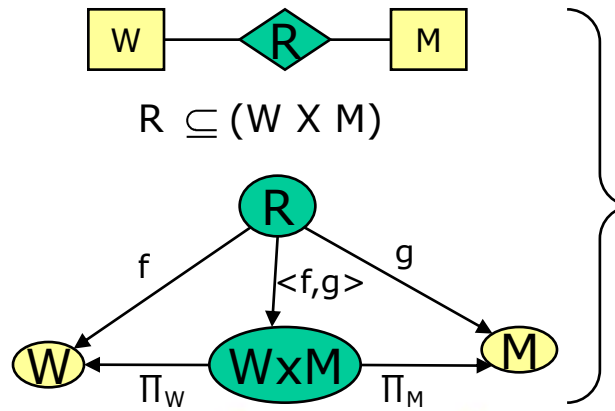
Jointly mono:
Binary relation
Ternary relation ...



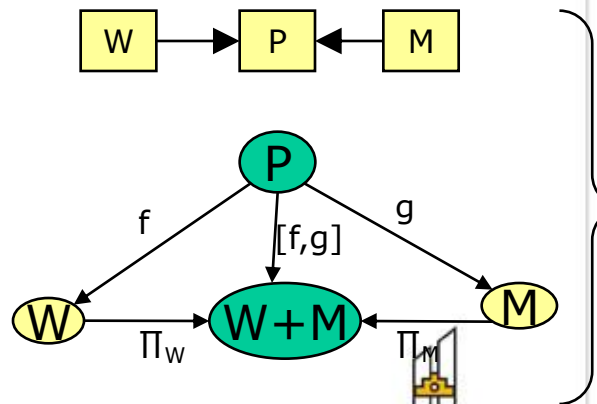
Disjoint union:
(inclusion)

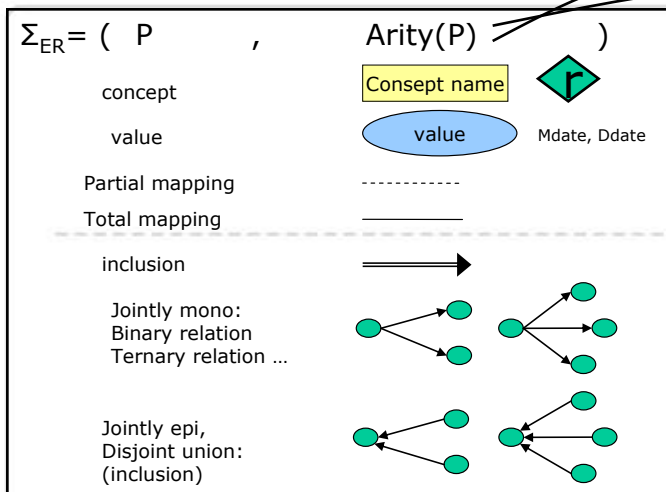


Equivalent ways for expressing binary relations. Categorically: product



Equivalent ways for expressing inclusion. Categorically: co-product





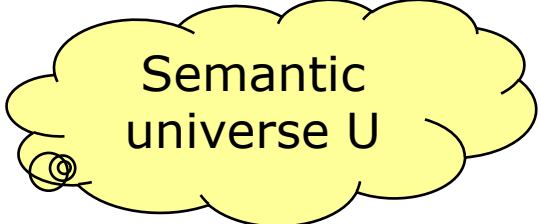
$$\Sigma_{ER} - sketch S = (G_S, D_S(P))$$

$$| \exists P \in \Sigma_{ER} : D_S(P) = \{d : Arity(P) \rightarrow G_S\}$$

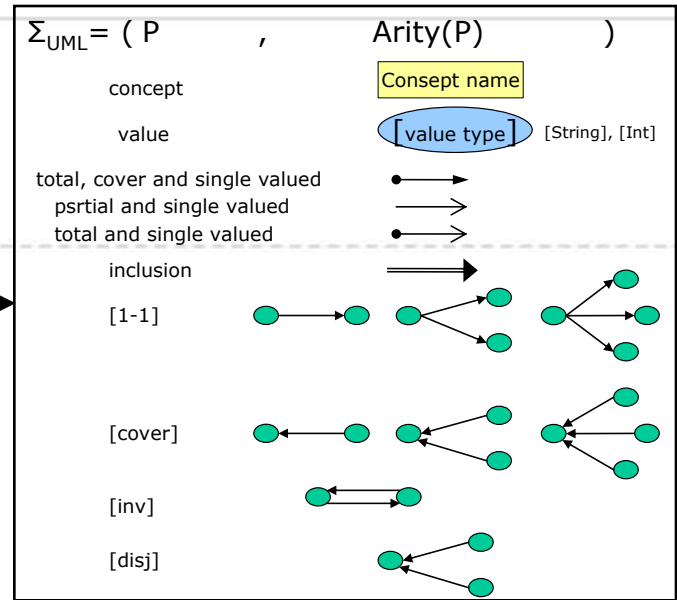
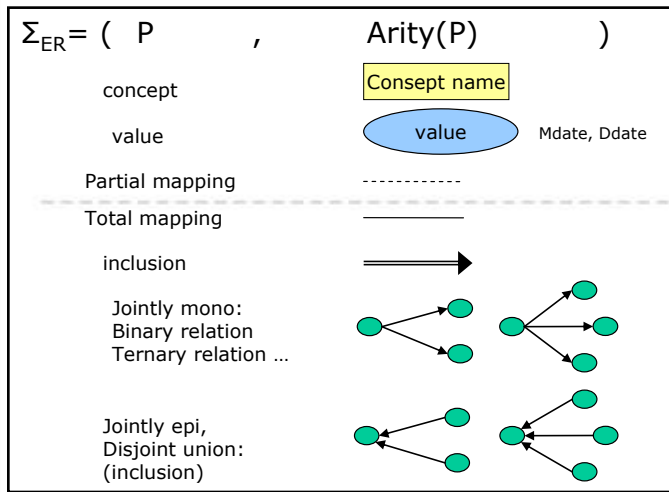
Semantic sketch of U

$$\Sigma_{ER} - sketch U = (G_U, D_U(P))$$

$$| \forall P \in \Sigma_{ER} : D_U(P) = \{\delta : Arity(P) \rightarrow G_U\}$$



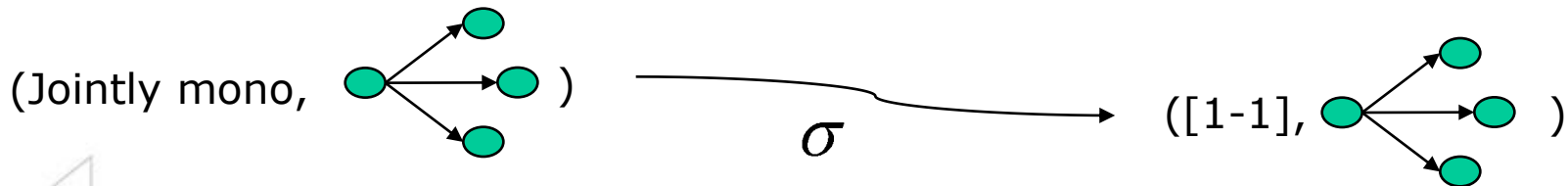
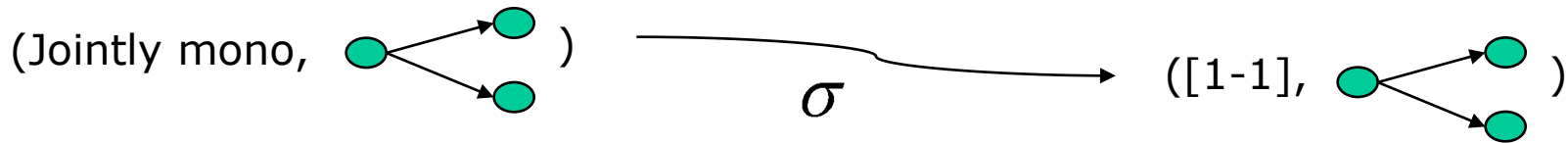
category SET (or graph U)



σ

Such that:

$$\forall P \in \Sigma_{ER} : Arity(P) \subseteq Arity(\sigma(P))$$



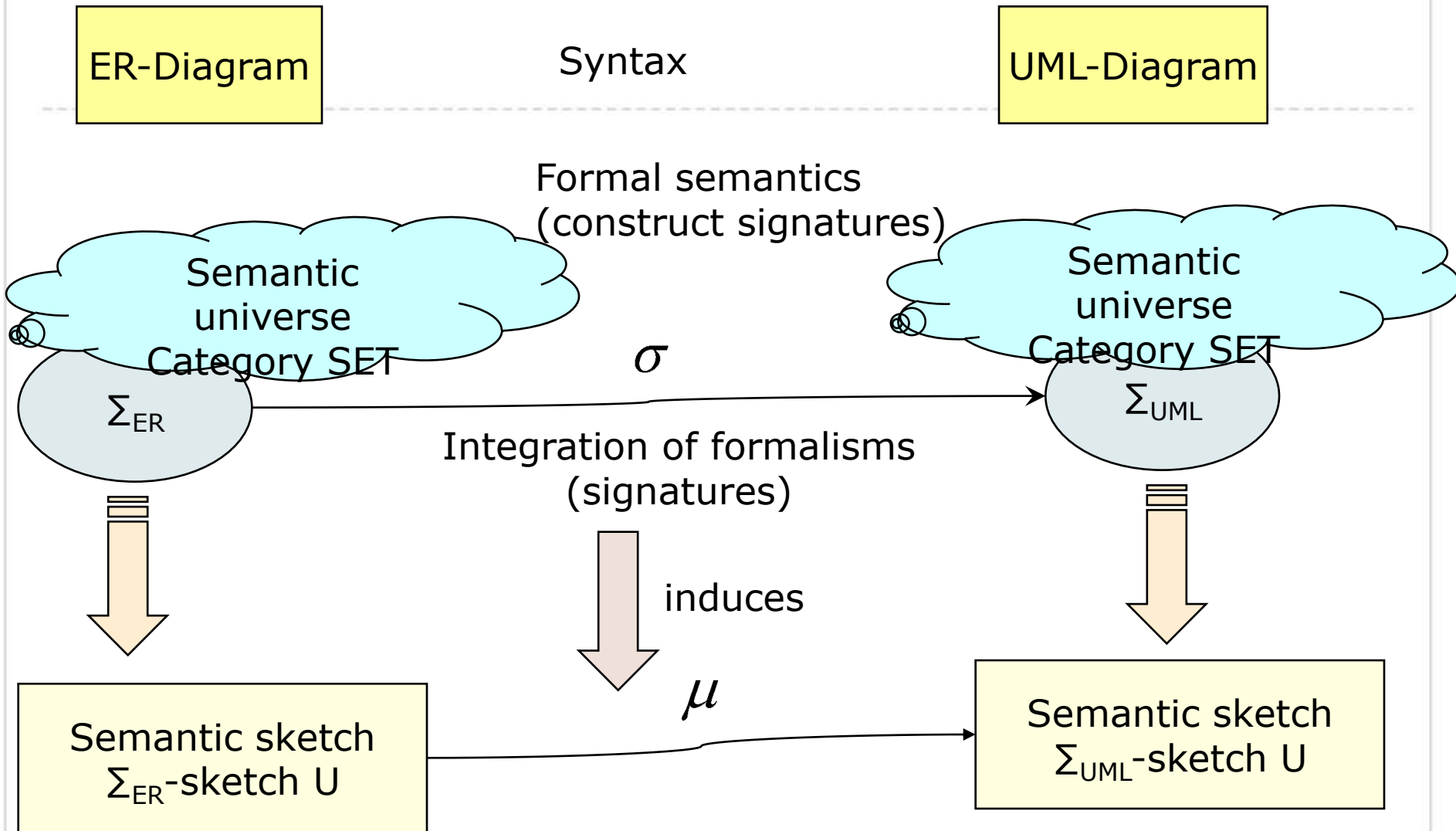
Our objectives (theory)

- Formalization of:
 - ER-diagrams,
 - UML-diagrams,
 - DB-schemes and
 - Ontology languages,by GS
- Investigation of the
 - integration,
 - combination and
 - modularizationof specifications within single specification formalism based on GS

Our objectives (technical)

- Design and development of tools supporting the application of GS in the field of software engineering:
 - drawing UML and ER diagrams based on GS formalism
 - support for mappings/translations between the two types of diagrams by mapping their signatures
 - code-generator
 - programming languages
 - specification language based on their graphical specifications
 - case studies to evaluate the theory in practice.

Summary



Thank you!

What's done and progress:

- Diskin's program
- Ørjan Hatlands Contribution
- New Master-thesis by Stian Skjerveggen

Facts

Sketch-able spec. \leftrightarrow Formal spec.

Diskin and Kadish:

- Graphical specifications (ER diagram or UML class diagram,) can be considered as:
 - abbreviations or
 - visualizations of sketches for a fixed signature where the signature is the corresponding diagram type

Diskin and Kadish:

- A graphical notation: a visualization on the top of a specification core
 - Visualization: a presentation (a user interface) of specification,
 - Specification "deals with" semantics of the notational constructs